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MashQL

A Data Mashup Language for the Data Web

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Motivation (Querying the Data Web)

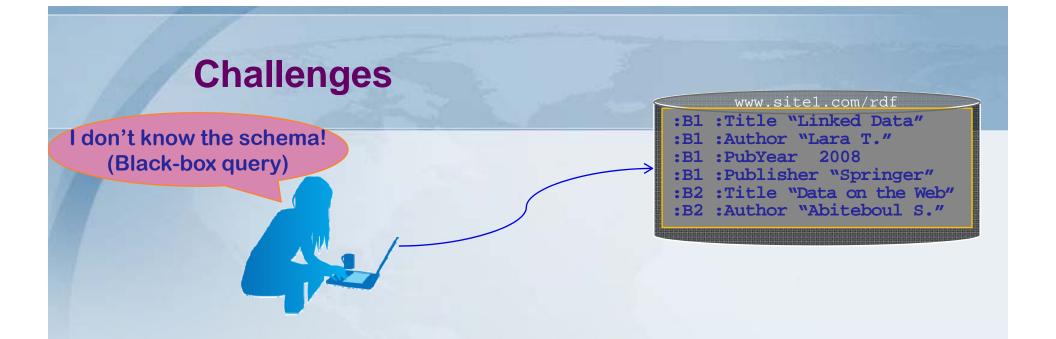
Can we query and mash up the Data Web <u>as simple as</u> filtering and piping Web Feeds?

→ We fundamentally investigate this problem from a Query Formulation viewpoint.

 \rightarrow A "data mashup" is a query.

Outline

- Challenges
- Related Work
- MashQL (A Query Formulation Language)
- Evaluation
- Conclusions and Discussion

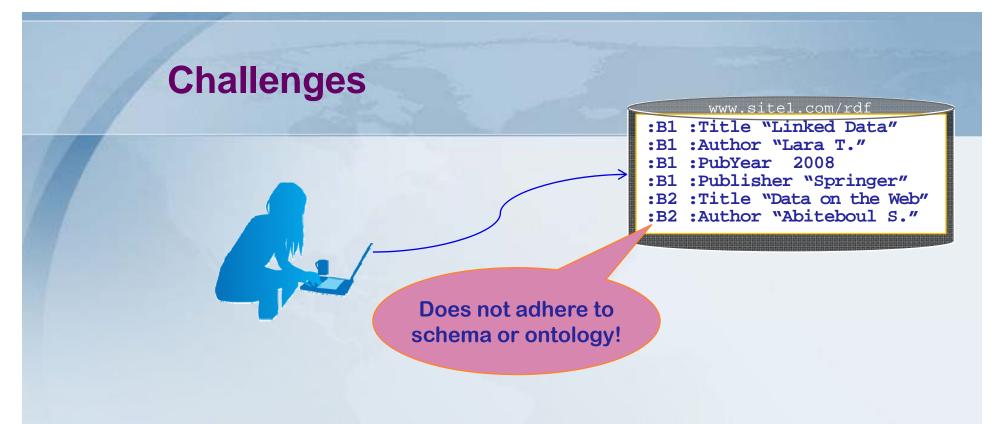


Problem Definition

We need a <u>way</u> to allow <u>end-users</u> formulate queries over structured data assuming that:

(1)

The user does not know the schema.



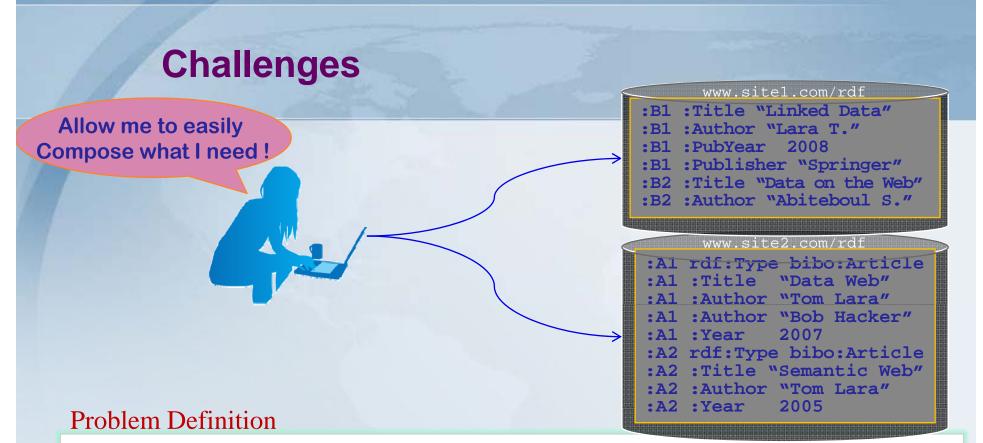
Problem Definition

How to allow *end-users* formulate queries over structured data assuming that:

(1)

(2)

- > The user does not know the schema.
- There is no offline or inline schema\ontology.



How to allow *end-users* formulate queries over structured data assuming that:

(1)

(2)

(3)

- > The user does not know the schema.
- There is no offline or inline schema\ontology.
- The query may involve multiple sources.



How to allow *end-users* formulate queries over structured data assuming that:

(1)

(2)

(3)

(4)

- The user does not know the schema.
- There is no offline or inline schema\ontology.
- The query may involve multiple sources.
- The query language is sufficiently expressive (not a single-purpose interface)

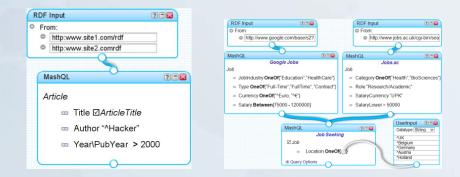
Related Work

Query formulation is the art of accessing and consuming structured data *easily*. (interdisciplinary subject)

Assumption	Query by Form	Query by Example	Conceptual Queries (ConQuer)	NL Queries	Interactive Queries (Lorel)	Visual Scripting (DeriPipes)
Don't know the schema	\checkmark	×	×		\checkmark	×
Schema-free data	×	×	×	×	-	\checkmark
Multiple sources	×	\checkmark	×	×	×	\checkmark
Expressive	×	\checkmark	\checkmark	\checkmark	-	\checkmark
Intuitive	\checkmark	×	×	\checkmark	\checkmark	×



A graphical query formulation Language.



all of these assumptions:

✓ The user does not know the schema.

Challenging combination

There is no offline or inline schema\ontology.
 The query may involve multiple sources.
 The query language is expressive.

(not a single-purpose interface)

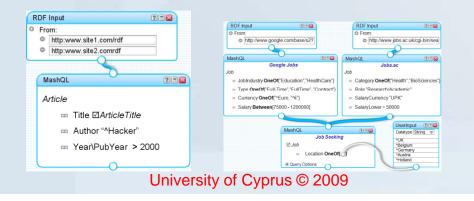
MashQL

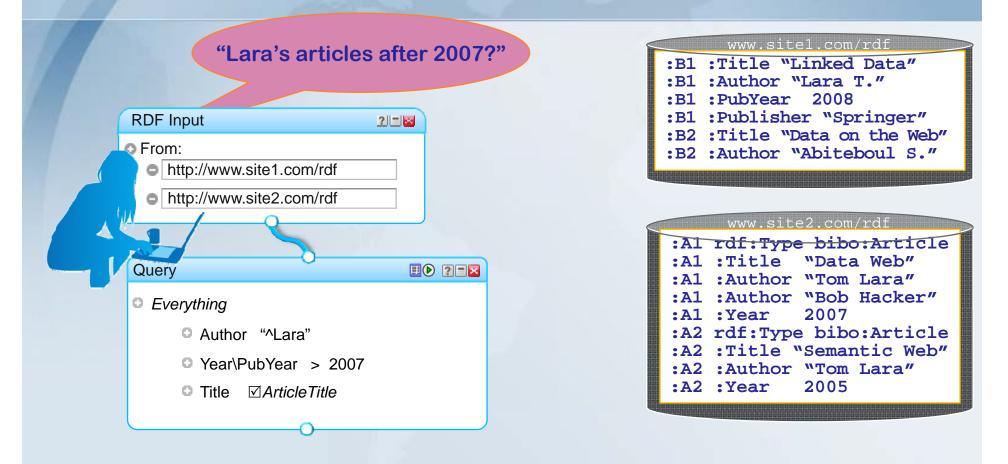
A general structured-data retrieval solution. (not merely an interface)

Without loosing this generality, we

Focus on RDF, as the most primitive query language, other data models can be mapped into it.

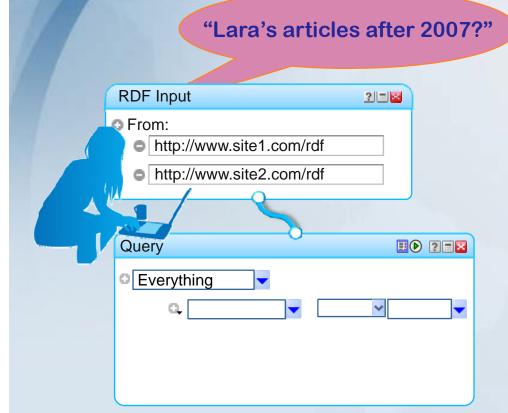
Follow Yahoo Pipes' visualization, in order to illustrate that Data Web can be queried and mashed up as Web Feeds.





Interactive Query Formulation.

MashQL queries are translated into and executed as SPARQL.



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	B1 :Title "Linked Data" B1 :Author "Lara T." B1 :PubYear 2008 B1 :Publisher "Springer" B2 :Title "Data on the Web" B2 :Author "Abiteboul S."
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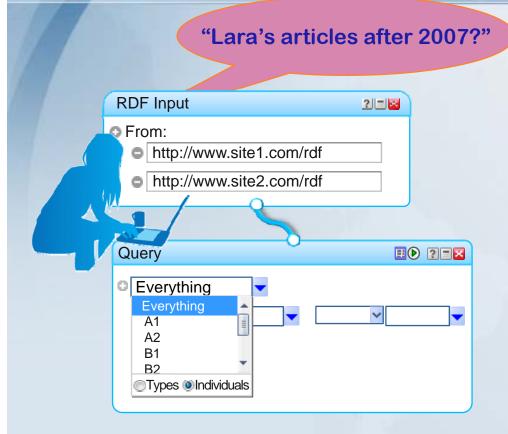
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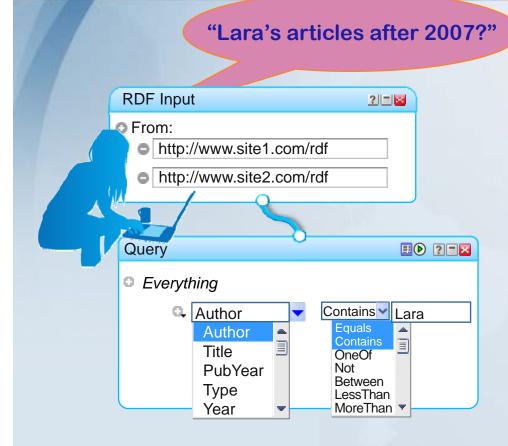


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:A1	:Author "Tom Lara"
:A1	:Author "Bob Hacker"
:A1	:Year 2007
:A2	rdf:Type bibo:Article
:A2	:Title "Semantic Web"
:A2	:Author "Tom Lara"
:A2	:Year 2005

Background queries

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SELECT X WHER	E {?S rdf:Type ?X
SELECT X WHER	E {?S ?P ?X}
SELECT X WHER	E {?X ?P ?O}



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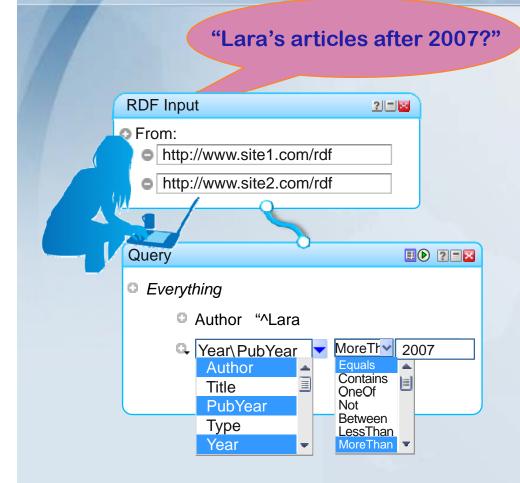
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:A2 :Author "Tom Lara"

:A2 :Year

Background query SELECT P WHERE {?Everything ?P ?0} University of Cyprus © 2009



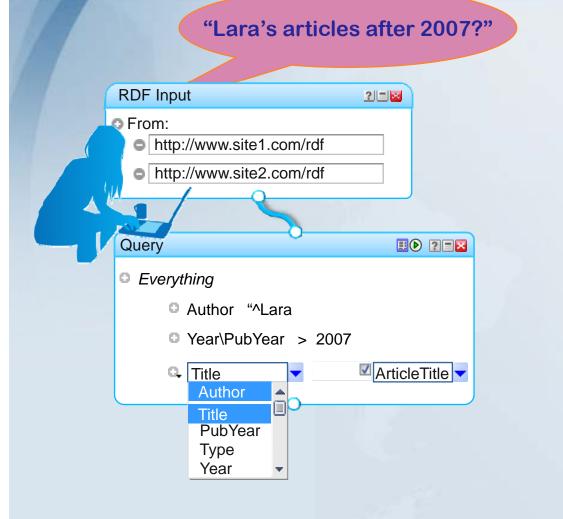
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:A2	:Author "Tom Lara"
:A2	:Year 2005

Background query

 SELECT P WHERE {?Everything ?P ?O}

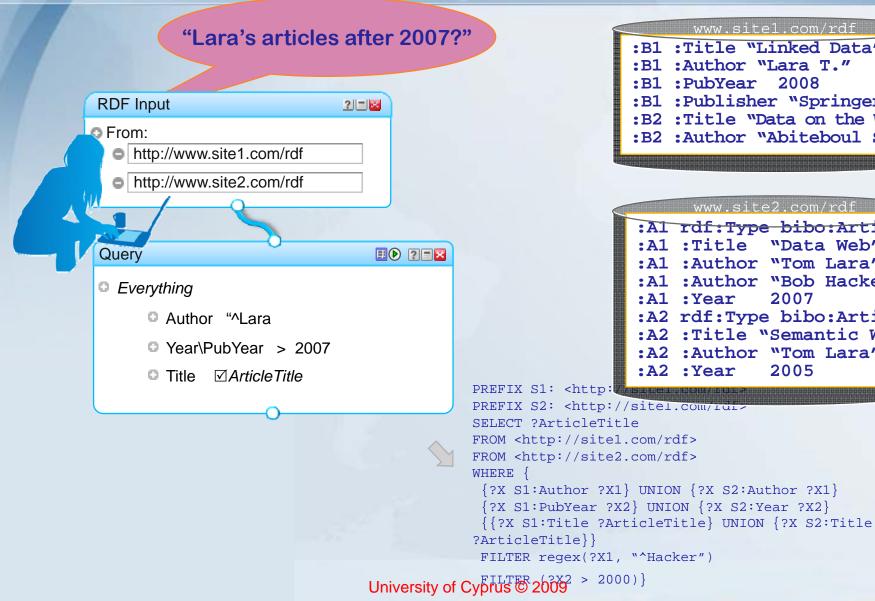
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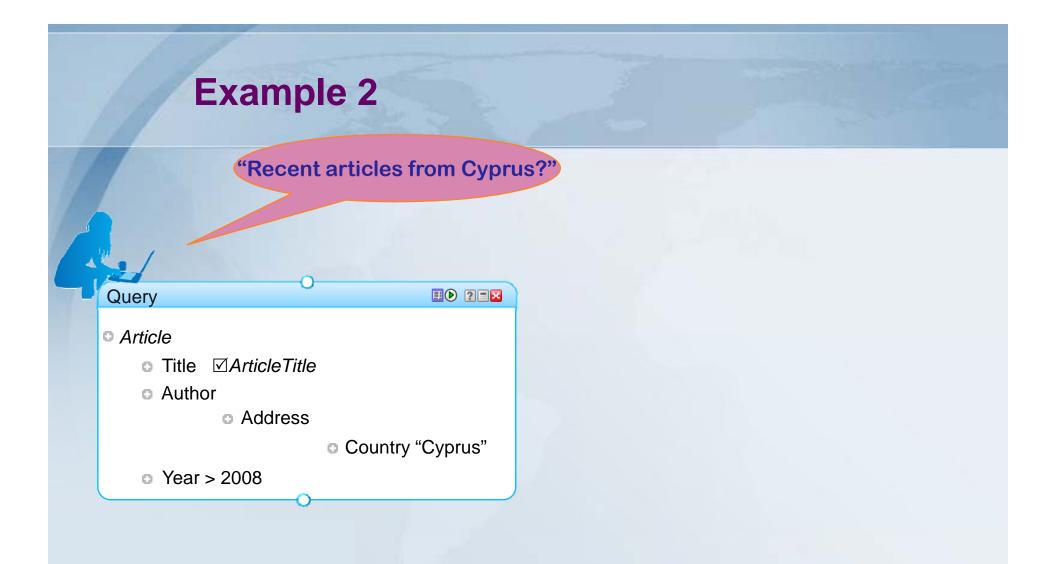
Background query SELECT P WHERE {?Everything ?P ?O} University of Cyprus © 2009



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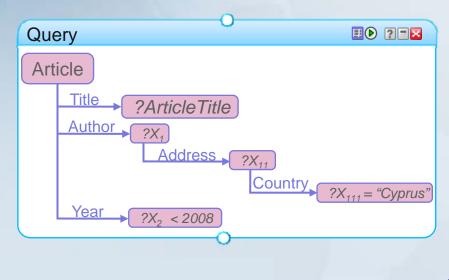
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	:A2	:Title "Semantic Web"
	:A2	:Author "Tom Lara"
	:A2	:Year 2005
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→ Retrieve every Article that: has a title, written by author, who has address, this address has a country called Cyprus, and the article is published after 2008.

The Intuition of MashQL

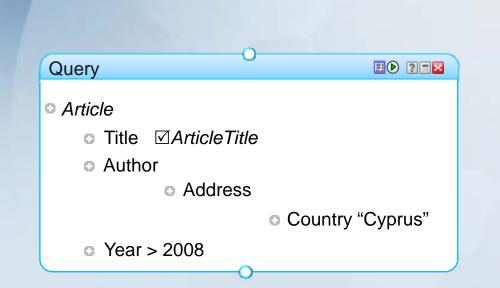


A query is a tree

- The root is called the query subject.
- Each branch is a *restriction*.
- Branches can be expanded, (information path)
- Object value filters

Def. A **Query** Q with a subject S, denoted by Q(S), is a set of restrictions on S. $Q(S) = R_1 AND ... AND R_n$. Dif. A **Subject** $S \in (I \cup V)$, where I is an identifier and V is a variable. Dif. A **Restriction** $R = \langle R_x, P, O_f \rangle$, where R_x is an optional restriction prefix that can be (maybe | without), P is a predicate ($P \in I \cup V$), and O_f is an object filter.

The Intuition of MashQL



An Object filter is one of :

- Equals
- Contains
- MoreThan
- LessThan
- Between
- one of
- Not(f)
- Information Path (sub query)

Def. An **object filter** $O_f = \langle O, f \rangle$, where O is an object and f is a filtering function one of :

 $O_f = \langle O \rangle$, where O is an object, $O \in V \cup I$.

 $O_f = \langle O, Equals(X, T, L_t) \rangle$, where X can be a variable or a constant, T is a datatype, and L_t is a language tag.

 $O_f = \langle O, Contains(X, T, L_t) \rangle$, where O is an object variable, X is a regex literal, T is a data type, and L_t is a language.

 $O_f = \langle O, MoreThan(X, T) \rangle$, where O is an object variable, X is a variable or a constant, T is a datatype.

 $O_f = \langle O, LessThan(X, T) \rangle$, where O is an object variable, X is a variable or a constant, T is a datatype identifier.

 $O_f = \langle O, Between(X, Y, T) \rangle$, where X and Y are variables or constants, T is a datatype identifier.

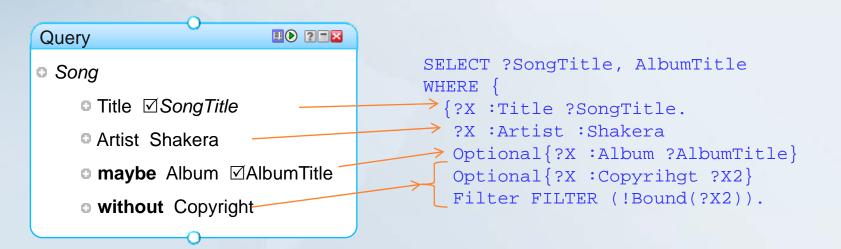
 $O_f = \langle O, OneOf(V) \rangle$, where O is an object variable, and V is a set of values $\{v_1, \dots, v_n\}$, v_i is a variable or constant.

 $O_f = \langle O, Not(f) \rangle$, where f is one of the functions defined above.

 $O_f = \langle O, Q_i(O) \rangle$, where O is an object ($O \in V \cup I$), and $Q_i(O)$ is a sub-query with O being the query subject.

More MashQL Constructs

Resection Operators {Required, Maybe, or Without} All restriction are required (i.e. AND), unless they are prefixed with "maybe" or "without"



More MashQL Constructs

Union operator (denoted as "\") between

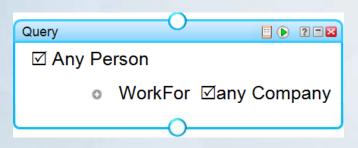
	MashQL	SELECT ?Person	
Objects	⊠ <i>Person</i>	WHERE { ?Person :WorkFor :Google UNION ?Person WorkFor :Yahoo}	
	MashQL 23	SELECT ?FName	
Predicates	Person □ Surname\Firstname ⊠FName	WHERE {	
Subjects	MashQL 2	<pre>SELECT ?AgentName, ?AgentPhone WHERE { {?Person rdf:type :Person.</pre>	
Subjects	a Person\Company	?Person :Name ?AgentName.	
	□ Name ⊠ <i>AgentName</i>	?Person :Phone ?AgentPhone}	
	□ Phone ⊠A <i>gentPone</i>	UNION {?Company rdf:type :Company.	
		<pre>?Company :Name ?AgentName. ?Company :Phone ?AgentPhone}}</pre>	
-	Query Dec 200	SELECT ?CustName,	
	Person	WHERE {	
Queries	o Name ⊠CustName ∖Company	?Person :Name ?CustName. UNION {?Company :Title ?CustName.	
	 Title	?Company :City ?X1.	

More MashQL Constructs

And several other constructs, including:

> Types

>



and Reverse Predicates



Datatypes and Language Tags

Formal Syntax and Semantics

Def.1 (Dataset): A dataset D is a set of triples, each triple t is formed as $\langle S, P, O \rangle$, where $S \in I, P \in I$, and $O \in I \cup L$.

Def.2 (Typed Literals): Every object literal must have a datatype D: If $O \in L$ then $O \in D$.

Def.3 (Language Tags): An object literal ($O \in L$) may have a language tag L_t

Def. 4 (Query): A Query Q with a subject S, denoted by Q(S), is a set of restrictions on S. $Q(S) = R_1 AND \dots AND R_n$.

Def. 5 (Subject): A subject $S \in (I \cup V)$, where I is an identifier and V is a variable.

Def. 6 (Restriction): A restriction $R = \langle R_x, P, O_f \rangle$, where R_x is an optional restriction prefix that can be (maybe | without), P is a predicate ($P \in I \cup V$), and O_f is an object.

Def.7 (Object Filter): An object filter $O_f = \langle O, f \rangle$, where O is an object and f is a filtering function. An object filter can have one of the following nine forms:

- 1. $O_f = \langle O \rangle$, where O is an object, $O \in V \cup I$. This is the simplest object filter, i.e., it does not add any restriction on the object value of the retrieved triples.
- 2. O_f = <O, Equals(X, T, L_i)>, where X can be a variable or a constant, T is a datatype, and L_t is a language tag. This filter restricts the retrieved results, such that, the object value O should be equal to X, with datatype T, and with language L_t.
- 3. _{Of} = <0, Contains(X, T, _{Lt})>, where O is an object variable, X is a regex literal, T is a data type, and _{Lt} is a language. This filter restricts the retrieved results, such that, the object value O should be equal to regex(X), with datatype T, and with language _{Lt}. A regex literal is a literal that contains a regular expression matching pattern.
- 4. O_i = <O, MoreThan(X, T)>, where O is an object variable, X is a variable or a constant, T is a datatype. This filter restricts the retrieved results, such that, the object value O should be more than X and with datatype T.
- 5. O_f = <O, LessThan(X, T)>, where O is an object variable, X is a variable or a constant, T is a datatype identifier. This filter restricts the retrieved results, such that, the object value O should be less than X and with datatype T (see rule-9).
- 6. O_f = <O, Between(X, Y, T)>, where X and Y are variables or constants, T is a datatype identifier. This filter restricts the retrieved results, such that, the object value O should be more than or equals X, less than or equals Y, and with datatype T.
- 7. $O_f = \langle O, OneOf(V) \rangle$, where O is an object variable, and V is a set of values { $v_1, ..., v_n$ }, v_i is a variable or constant. This filter restricts the retrieved results, such that, the object value O should be equal to one of the values in V.
- 8. O_f = <O, Not(f)>, where f is one of the functions defined above. This filter extends all of the above functions with simple negation. The filter is same as the *Equals* filter but with negation, i.e., Not Equal.
- 9. $O_f = \langle O, Q_i(O) \rangle$, where O is an object ($O \in V \cup I$), and $Q_i(O)$ is a sub-query with O being the query subject. The restrictions defined in the sub-query $Q_i(O)$ should be satisfied as well. Notice that this definition is recursive; however, this does not mean the query itself is recursive.

Def.8 (Types): A subject ($S \in I$) or an object ($O \in I$) can be prefixed with "a" or "an" to mean the instances of this subject/object type, instead of the subject/object itself.

Def.9 (Union): A union can be declared between objects, predicates, subjects and/or queries, in the following forms:

- 1. $O_n = \langle O_1 | O_2 | \dots | O_n \rangle$, to indicate unions between objects, where $O_i \in I$.
- 2. $P_n = \langle P_1 | P_2 | \dots | P_n \rangle$, to indicate unions between predicates, where $P_i \in I$.
- 3. $S_n = \langle S_1 | S_2 | \dots | S_n \rangle$, to indicate unions between subjects, where $S_i \in I$.

4. $Q_n = \langle Q_1 | Q_2 | \dots | Q_n \rangle$, to indicate unions between queries.

Def.10 (Reverse): $\langle P \rangle$ indicates the reverse of the predicate P. Let R_1 be a restriction on S such that $\langle S P O \rangle$, and R_2 be $\langle O P S \rangle$, R_1 and R_2 have the same meaning.

Query Formulation Algorithm

Formalization of the background queries

Select Subject S

(1) $S \in S_T : \pi_O(\sigma_{P=:Type'}(D))$ (2) $S \in S_I : \pi_S(D) \cup \pi_O(\sigma_{O\in}(D))$ (3) $S \in V$ (1) $O1: \{(?S1 <:Type> ?O1)\}$ (1) $O1: \{(?S1 ?P1 ?O1)\}$ UNION $O1: \{(?S1 ?P1? O1)\}$ (1) $O1: \{(?S1 ?P1 ?O1)\}$ (1) $O1: \{(?S1 ?P1? O1)\}$ (2) $O1: \{(?S1 ?P1? O1)\}$ (2) $O1: \{(?S1 ?P1? O1)\}$ (3) $S \in V$

Select a property P

(4) $(S \in S_T) \rightarrow P \in \pi_{P2} (\sigma_{P1=:Type' \land O1=Subject}(D) \rtimes_{S1=S2} \sigma(D))$ (5) $(S \in S_I) \rightarrow P \in \pi_P (\sigma_{S=Subject}(D))$ (6) $(S \in V) \rightarrow P \in \pi_P (\sigma(D))$ (7) $P \in V$ (4') P2:{(?S1 <:Type> <S>)(?S2 ?P2 ?O2)}
(5') P1:{(<S> ?P1 ?O1)}
(6') P1:{(?S1 ?P1 ?O1)}

```
Add a filter on P

(8) (S \in S_I) \land (P \in V) \rightarrow O \in \pi_{OI}(\sigma_{SI=S \land OI \in}(D))

(9) (S \in S_I) \land (P \notin V) \rightarrow O \in \pi_{OI}(\sigma_{SI=S \land PI=P \land OI \in}(D))

(10) (S \in S_T) \land (P \in V) \rightarrow O \in \pi_{O2}(\sigma_{PI=':Type' \land OI=S}(D) \rtimes_{SI=S2} \sigma(D))

(11) (S \in S_T) \land (P \notin V) \rightarrow O \in \pi_{O2}(\sigma_{PI=':Type' \land OI=S}(D) \rtimes_{SI=S2} \sigma_{P2=P})

(11) (S \in V) \land (P \notin V) \rightarrow O \in \pi_{O2}(\sigma_{PI=':Type' \land OI=S}(D) \rtimes_{SI=S2} \sigma_{P2=P})

(11') O: \{(?S < rdf:Type > <S >)(?S < P > ?O)\}

(12) (S \in V) \land (P \notin V) \rightarrow O \in \pi_{O}(\sigma(D))

(13) (S \in V) \land (P \notin V) \rightarrow O \in \pi_{O}(\sigma_{P=P}(D))

(14') O: \{(?S1 < P1 ?O1)\}

(13') O: \{(?S1 < P > ?O1)\}
```

MashQL-SPARQL Mapping Rules

Rule-1: The symbol \square before a variable means that it will be returned in the results; i.e., included in the SELECT part of in SPARQL. If the output of the query is input to another, use "CONSTRUCT *".

Rule-2: In any of the following rules, if a subject, predicate, or object is italicized: it is seen as a SPARQL variable, i.e. prefixed with "?".

Also mapped into

Oracle's SPARQL

Rule-3: If S is a subject and $R = \langle P, Of \rangle$, the mapping is: $\{S P O\}$.

Rule-4: If S is a subject and R = <maybe, P, Of>, the mapping is: {OPTIONAL{S P O}}.

Rule-5: If S is a subject and R = < without, P, Of>, the mapping is: {S P O. FILTER (!bound(?O))}.

Rule 6. If Of = <O, Equals(X, T, Lt)>:

Append the mapping with: FILTER(?O = X)

If T ≠ Null: Append the mapping with: FILTER(datatype(?O)=T)

If Lt ≠ Null: Append the mapping with: FILTER(lang(?O) = Lt)

Rule 7. If Of = Cor

Append the ma

If T ≠ Null: App

If Lt ≠ Null: App

Rule 8. If Of = <0,

Append the ma

If T ≠ Null: App

Rule 9. If Of = <0, 2000 man(x, 1)*. Append the mapping with: FILTER(?O < X)

If $T \neq Null$: Append the mapping with: FILTER(datatype(?O=T))

Rule 10. If Of = <0, Between(X, Y, T)>:

Append the mapping with: FILTER(?O >=X)&& FILTER(?O<=Y)

```
If T ≠ Null: Append the mapping with: FILTER(datatype(?O)=T)
```

Rule 11. If Of = <0, OneOf (V)>:

Append the mapping with: {FILTER(?O = V1)|| . . . || FILTER(?O = Vn)}

If Vi is a regex-ed literal, the ith filter above should be replaced with: FILTER Regex(?O, Vi)

Rule 12. If Of = <O, Not(f)>: The f filter will be generated as above, but with a negation.

Rule 13. If $Of = \langle O, Qi(O) \rangle$: Repeat all mapping rules to generate Qi(O).

Rule 14. If a subject S is prefixed with "a" or "an": Append the mapping with: {?S rdf:type :S}

Rule 15. If an object O is prefixed with "a" or "an": Append the mapping with: {?O rdf:type :O}

Rule 16. Given On , If n >1 and Oi \in I : The mapping in rules 3-4 will be:{{S P :O1} UNION . . . UNION {S P :On}}

Rule 17. Given Pn , If n >1 and Pi \in I : The mapping in rules 3-4 will be: {{S :P1 O} UNION . . . UNION {S :Pn O}}

Rule 18. Given Sn , If n >1 and Si \in I : Regenerate the query n times, each time with Si as a root, and with a UNION between the queries.

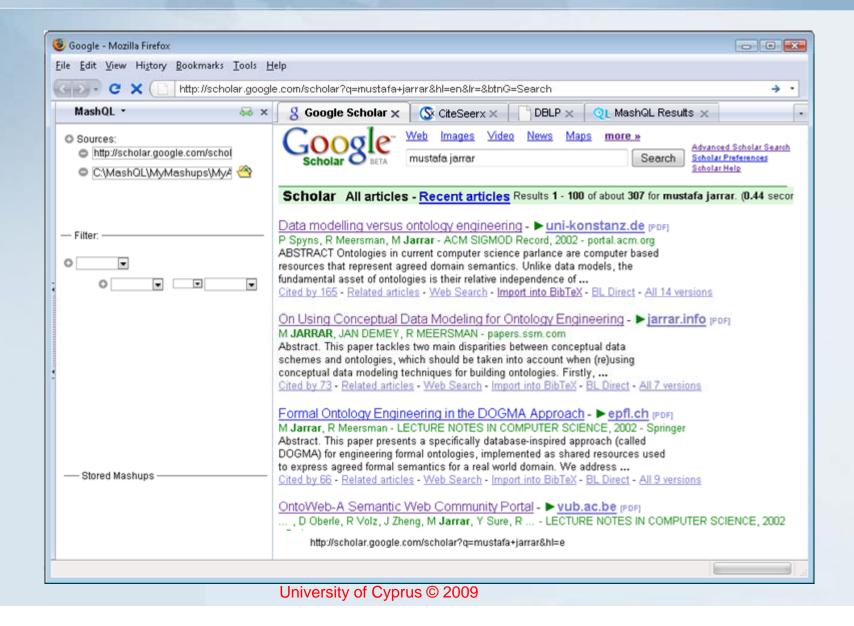
Rule 19. Given Qn, If n > 1: Add UNION between the n queries.

Rule 20. If S is a subject and $R = \langle P, O \rangle$, the mapping is: {O P S}.

MashQL Editor

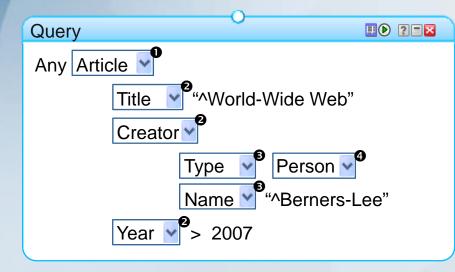
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(but some URIs are too cryptic)							
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MashQL Firefox Add-On (Light-mashups @ your browser)



Evaluation (DBLP, Experiment 1)

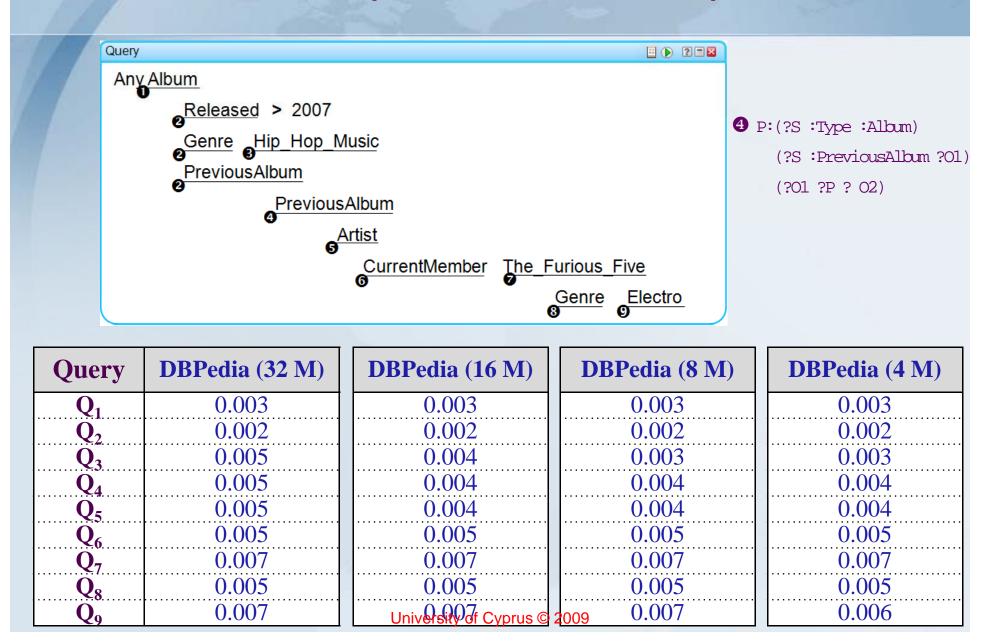
(User Interaction Response Time) How long it takes to generate the next list?



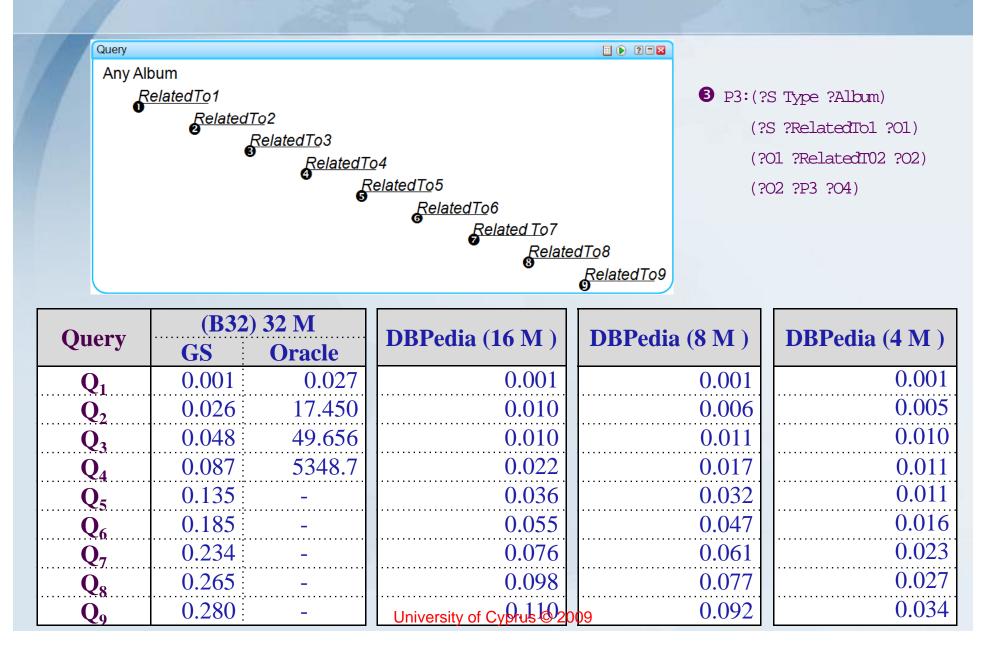
- **0** 0:(?S Type ?0)
- 2 P:(?S Type Article)(?S ?P ?01)
- 3 P:(?S Type Article) (?S Creator ?01) (?01 ?P ?02)
- ④ 0:(?S Type Article)
 (?S Creator ?01)(?01 Type ?0)

Ouony	DBLP (9M triples)		DBLP (4M triples)		DBLP (2M triples)	
Query	GS	Oracle	GS	Oracle	GS	Oracle
Q_1	0.003	0.005	0.003	0.004	0.003	0.003
Q ₂	0.001	0.136	0.001	0.148	0.001	0.108
Q ₃	0.001	0.187	0.001	0.846	0.001	0.671
Q ₄	0.001	1.208	0.001	0.835	0.001	0.650

Evaluation (DBPedia, Experiment 2)



Evaluation (DBPedia, Experiment 3)



Conclusions

"Query and mash up the Data Web as simple as filtering up Web Feeds" is a query formulation problem.

- End-users can navigate, query, and mash up unknown graphs.
 without knowing the schema. Data is schema-free. Multiple sources.
- MashQL is expressive as SPARQL. Except NAMED GRAPH.
- MashQL is not merely a SPARQL interface, or limited RDF.
 It has its own path-pattern intuition (can be similarly used for XML and DB).

Future Work



- Reasoning, Keyword Search, Aggregation Functions, etc.
- Results Presentation (Should be tacked fundamentally).
- Firefox add-on Mashup/query editor.

RDF summaries for SPARQL optimization.

Thank You

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